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GLENN PATENT GROUP
3475 Edison Way
Suite L
Menlo Park, CA 94025

EXAMINER	
HENNING, MATTHEW T	

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Please find below and/or attached an Office communication concerning this application or proceeding.

The time period for reply, if any, is set in the attached communication.

Office Action Summary	**Application No.** 09/913,686	**Applicant(s)** RUMP ET AL.	
	Examiner Matthew T. Henning	**Art Unit** 2131	

-- The MAILING DATE of this communication appears on the cover sheet with the correspondence address --

Period for Reply

A SHORTENED STATUTORY PERIOD FOR REPLY IS SET TO EXPIRE 3 MONTH(S) OR THIRTY (30) DAYS, WHICHEVER IS LONGER, FROM THE MAILING DATE OF THIS COMMUNICATION.

- Extensions of time may be available under the provisions of 37 CFR 1.136(a). In no event, however, may a reply be timely filed after SIX (6) MONTHS from the mailing date of this communication.
- If NO period for reply is specified above, the maximum statutory period will apply and will expire SIX (6) MONTHS from the mailing date of this communication.
- Failure to reply within the set or extended period for reply will, by statute, cause the application to become ABANDONED (35 U.S.C. § 133). Any reply received by the Office later than three months after the mailing date of this communication, even if timely filed, may reduce any earned patent term adjustment. See 37 CFR 1.704(b).

Status

1)☒ Responsive to communication(s) filed on *06 June 2006*.

2a)□ This action is **FINAL**. 2b)☒ This action is non-final.

3)□ Since this application is in condition for allowance except for formal matters, prosecution as to the merits is closed in accordance with the practice under *Ex parte Quayle*, 1935 C.D. 11, 453 O.G. 213.

Disposition of Claims

4)☒ Claim(s) *1-30* is/are pending in the application.

4a) Of the above claim(s) _____ is/are withdrawn from consideration.

5)□ Claim(s) _____ is/are allowed.

6)☒ Claim(s) *1-30* is/are rejected.

7)□ Claim(s) _____ is/are objected to.

8)□ Claim(s) _____ are subject to restriction and/or election requirement.

Application Papers

9)□ The specification is objected to by the Examiner.

10)☒ The drawing(s) filed on *24 January 2002* is/are: a)☒ accepted or b)□ objected to by the Examiner.

Applicant may not request that any objection to the drawing(s) be held in abeyance. See 37 CFR 1.85(a).

Replacement drawing sheet(s) including the correction is required if the drawing(s) is objected to. See 37 CFR 1.121(d).

11)□ The oath or declaration is objected to by the Examiner. Note the attached Office Action or form PTO-152.

Priority under 35 U.S.C. § 119

12)☒ Acknowledgment is made of a claim for foreign priority under 35 U.S.C. § 119(a)-(d) or (f).

a)☒ All b)□ Some * c)□ None of:

1.□ Certified copies of the priority documents have been received.

2.□ Certified copies of the priority documents have been received in Application No. _____.

3.☒ Copies of the certified copies of the priority documents have been received in this National Stage application from the International Bureau (PCT Rule 17.2(a)).

* See the attached detailed Office action for a list of the certified copies not received.

Attachment(s)

1) ☒ Notice of References Cited (PTO-892)
2) □ Notice of Draftsperson's Patent Drawing Review (PTO-948)
3) □ Information Disclosure Statement(s) (PTO/SB/08) Paper No(s)/Mail Date _____.
4) □ Interview Summary (PTO-413) Paper No(s)/Mail Date. _____ .
5) □ Notice of Informal Patent Application
6) □ Other: _____.

This action is in response to the communication filed on 6/6/2007.

DETAILED ACTION

Continued Examination Under 37 CFR 1.114

A request for continued examination under 37 CFR 1.114, including the fee set forth in 37 CFR 1.17(e), was filed in this application after final rejection. Since this application is eligible for continued examination under 37 CFR 1.114, and the fee set forth in 37 CFR 1.17(e) has been timely paid, the finality of the previous Office action has been withdrawn pursuant to 37 CFR 1.114. Applicant's submission filed on 6/6/2007 has been entered.

Response to Arguments

Applicant's arguments filed 6/6/2007 have been fully considered but moot in view of the new grounds of rejection presented below.

Claims 1-30 have been examined and claim 31 has been cancelled.

All objections and rejections not set forth below have been withdrawn.

Claim Rejections - 35 USC § 103

The following is a quotation of 35 U.S.C. 103(a) which forms the basis for all obviousness rejections set forth in this Office action:

> (a) A patent may not be obtained though the invention is not identically disclosed or described as set forth in section 102 of this title, if the differences between the subject matter sought to be patented and the prior art are such that the subject matter as a whole would have been obvious at the time the invention was made to a person having ordinary skill in the art to which said subject matter pertains. Patentability shall not be negatived by the manner in which the invention was made.

Claims 1-7, 14, 16-17, 19, 23, 25-29 are rejected under 35 U.S.C. 103(a) as being unpatentable over Van Oorschot et al. (US Patent Number 5,850,443) hereinafter referred to as

Van Oorschot, and further in view of Nardone et al. (US Patent Number 5,805,700) hereinafter referred to as Nardone, and further in view of Yatsukawa (US Patent Number 6,148,404).

Regarding claim 1, Van Oorschot disclosed a method for producing a payload data stream comprising a header and a payload data block containing encrypted payload data (See Van Oorschot Fig. 3 X-fields, header fields, and encrypted message field), comprising the following steps: generating a payload data key for a payload data encryption algorithm for encrypting payload data (See Van Oorschot Col. 6 Lines 41-43 and Fig. 3 "Create low trust symmetric key" K'); encrypting a first section of the payload data using said payload data key and said payload data encryption algorithm to obtain an encrypted section of said payload data block of said payload data stream (See Van Oorschot Col. 6 Lines 42-43 and Fig. 3 "Symmetric encryption" and "encrypted message"), said first section including audio data, video data, a combination of audio data and video data, text data, or binary data forming an executable program (See Van Oorschot Abstract ciphertext), wherein a second section of the payload data remains unencrypted (See Van Oorschot Col. 6 Lines 45-47 "public key of entity A"); processing the unencrypted section of said payload data (See Van Oorschot Col. 6 Lines 45-50 "hash of X" which contains the public key of A) to deduce information characterizing the unencrypted second section of said payload data (See Van Oorschot Col. 6 Lines 49-60 h40(X)); linking said information and said payload data key by means of an invertible logic linkage to obtain a basic value (See Van Oorschot Col. 6 Lines 56-60 "K' XOR h40(X)"); encrypting said basic value using a key of two keys being different from each other by an asymmetrical encryption method, said two different keys being the public and the private keys respectively for said asymmetrical encryption method, to obtain an output value being an encrypted version of said payload data key (See Van Oorschot Col. 6 Line 60 – Col. 7 Line 7); and entering said output value into said header of said payload data stream (See Van Oorschot Col. 6 Line 65 – Col. 7 Line 7 and Fig. 3 "A's header field" and "B's header field"), but Van Oorschot failed to disclose that the second section included audio data, video data, a combination of audio data and video data, text data, or binary data forming an executable program, or the X-fields containing audio data, video data, a combination of audio data and video data, text data or binary data forming an executable program.

Nardone teaches that movie data needs to be protected from being copied and that this is generally done through encrypting the movie data (See Nardone Col. 1 Lines 22-37), and further that in order to save on processing cost, only portions of the movie data should be encrypted (See Nardone Col. 1 Summary of the Invention).

Yatsukawa teaches that by providing a public key in a digital certificate, the public key can be authenticated by the recipient of the key (See Yatsukawa Col. 5 Paragraph 1).

It would have been obvious to the ordinary person skilled in the art at the time of invention to employ the teachings of Nardone in the encryption system of Van Oorschot by encrypting video data, and further by only encrypting portions of the data. This would have been obvious because the ordinary person skilled in the art would have been motivated to protect movie data and to save on processing cost. It further would have been obvious to the ordinary person skilled in the art at the time of invention to employ the teachings of Yatsukawa in the encryption system of Van Oorschot by providing the public key in a public key certificate in the X-fields. This would have been obvious because the ordinary person skilled in the art would

have been motivated to provide the recipient with means to verify the authenticity of the public key. Furthermore, it was well know that public key certificates contain text data.

Regarding claim 17, Van Oorschot disclosed a method for decrypting an encrypted payload data stream comprising a header and a payload data block containing a first section having encrypted payload data (encrypted message), said first section including audio data, video data, a combination of audio data and video data, text data, or binary data forming an executable program (See Van Oorschot Abstract ciphertext), and a second section having unencrypted payload data (public key of A), said header comprising an output value having been generated by an encryption of a basic value by an asymmetrical encryption method using a key of two different keys including a private and a public key, said basic value representing a linkage of a payload data key, with which said first section having encrypted payload data is encrypted using a payload data encryption algorithm, and information deduced by a certain processing of the unencrypted second section of the payload data, said information characterizing a certain part of said payload data stream unambiguously (See rejection of claim 1 above), said method comprising the following steps: obtaining said output value from said header (See Van Oorschot Fig. 4 “B’s Header Field” and Col. 4 Lines 51-52); decrypting said output value using the other key of said asymmetrical encryption method to obtain said basic value (See Van Oorschot Fig. 4 “private key decryption” and “”B’s high trust private key” and Col. 4 Lines 53-54); processing the unencrypted second section of said payload data stream using the processing method used when encrypting to deduce information characterizing the unencrypted second (See Van Oorschot Fig. 4 “X-fields” and Col. 6 Lines 45-47); linking said information and said basic value using the corresponding linkage as it has been used when encrypting to obtain said payload data key (See Van Oorschot Fig. 4 “Unlevelling” and “X-fields” and Col. 4 Lines 54-56); and decrypting the first section containing the encrypted payload data using said payload data key and said payload data encryption algorithm used when encrypting (See Van Oorschot Fig. 4 “symmetric decryption” and “message”), but Van Oorschot failed to disclose that the second section included audio data, video data, a combination of audio data and video data, text data, or binary data forming an executable program, or the X-fields containing audio data, video data, a combination of audio data and video data, text data or binary data forming an executable program.

Nardone teaches that movie data needs to be protected from being copied and that this is generally done through encrypting the movie data (See Nardone Col. 1 Lines 22-37), and further that in order to save on processing cost, only portions of the movie data should be encrypted (See Nardone Col. 1 Summary of the Invention).

Yatsukawa teaches that by providing a public key in a digital certificate, the public key can be authenticated by the recipient of the key (See Yatsukawa Col. 5 Paragraph 1).

It would have been obvious to the ordinary person skilled in the art at the time of invention to employ the teachings of Nardone in the encryption system of Van Oorschot by encrypting video data, and further by only encrypting portions of the data. This would have been obvious because the ordinary person skilled in the art would have been motivated to protect movie data and to save on processing cost. It further would have been obvious to the ordinary person skilled in the art at the time of invention to employ the teachings of Yatsukawa in the encryption system of Van Oorschot by providing the public key in a public key certificate in the X-fields. This would have been obvious because the ordinary person skilled in the art would

have been motivated to provide the recipient with means to verify the authenticity of the public key. Furthermore, it was well know that public key certificates contain text data.

Regarding claim 28, Van Oorschot disclosed a device for producing a payload data stream comprising a header and a payload data block containing encrypted payload data (See Van Oorschot Fig. 3 X-fields, header fields, and encrypted message field), comprising: a generator for generating a payload data key for a payload data encryption algorithm for encrypting payload data (See Van Oorschot Col. 6 Lines 41-43 and Fig. 3 "Create low trust symmetric key" K'); a first encryptor for encrypting a first section of the payload data using said payload data key and said payload data encryption algorithm to obtain an encrypted section of said payload data block of said payload data stream (See Van Oorschot Col. 6 Lines 42-43 and Fig. 3 "Symmetric encryption" and "encrypted message"), said first section including audio data, video data, a combination of audio data and video data, text data, or binary data forming an executable program (See Van Oorschot Abstract ciphertext), wherein a second section of the payload data remains unencrypted (See Van Oorschot Col. 6 Lines 45-47 "public key of entity A"); a processor for processing the unencrypted section of said payload data (See Van Oorschot Col. 6 Lines 45-50 "hash of X" which contains the public key of A) to deduce information characterizing the unencrypted second section of said payload data (See Van Oorschot Col. 6 Lines 49-60 h40(X)); a linker for linking said information and said payload data key by means of an invertible logic linkage to obtain a basic value (See Van Oorschot Col. 6 Lines 56-60 "K' XOR h40(X)"); a second encryptor for encrypting said basic value using a key of two keys being different from each other by an asymmetrical encryption method, said two different keys being the public and the private keys respectively for said asymmetrical encryption method, to obtain an output value being an encrypted version of said payload data key (See Van Oorschot Col. 6 Line 60 – Col. 7 Line 7); and entering said output value into said header of said payload data stream (See Van Oorschot Col. 6 Line 65 – Col. 7 Line 7 and Fig. 3 "A's header field" and "B's header field"), but Van Oorschot failed to disclose that the second section included audio data, video data, a combination of audio data and video data, text data, or binary data forming an executable program, or the X-fields containing audio data, video data, a combination of audio data and video data, text data or binary data forming an executable program.

Nardone teaches that movie data needs to be protected from being copied and that this is generally done through encrypting the movie data (See Nardone Col. 1 Lines 22-37), and further that in order to save on processing cost, only portions of the movie data should be encrypted (See Nardone Col. 1 Summary of the Invention).

Yatsukawa teaches that by providing a public key in a digital certificate, the public key can be authenticated by the recipient of the key (See Yatsukawa Col. 5 Paragraph 1).

It would have been obvious to the ordinary person skilled in the art at the time of invention to employ the teachings of Nardone in the encryption system of Van Oorschot by encrypting video data, and further by only encrypting portions of the data. This would have been obvious because the ordinary person skilled in the art would have been motivated to protect movie data and to save on processing cost. It further would have been obvious to the ordinary person skilled in the art at the time of invention to employ the teachings of Yatsukawa in the encryption system of Van Oorschot by providing the public key in a public key certificate in the X-fields. This would have been obvious because the ordinary person skilled in the art would

have been motivated to provide the recipient with means to verify the authenticity of the public key. Furthermore, it was well know that public key certificates contain text data.

Regarding claim 29, Van Oorschot disclosed a device for decrypting an encrypted payload data stream comprising a header and a payload data block containing a first section having encrypted payload data (encrypted message), said first section including audio data, video data, a combination of audio data and video data, text data, or binary data forming an executable program (See Van Oorschot Abstract ciphertext), and a second section having unencrypted payload data (public key of A), said header comprising an output value having been generated by an encryption of a basic value by an asymmetrical encryption method using a key of two different keys including a private and a public key, said basic value representing a linkage of a payload data key, with which said first section having encrypted payload data is encrypted using a payload data encryption algorithm, and information deduced by a certain processing of the unencrypted second section of the payload data, said information characterizing a certain part of said payload data stream unambiguously (See rejection of claim 1 above), said device further comprising: means for obtaining said output value from said header (See Van Oorschot Fig. 4 “B’s Header Field” and Col. 4 Lines 51-52); a first decryptor for decrypting said output value using the other key of said asymmetrical encryption method to obtain said basic value (See Van Oorschot Fig. 4 “private key decryption” and “”B’s high trust private key” and Col. 4 Lines 53-54); a processor for processing the unencrypted second section of said payload data stream using the processing method used when encrypting to deduce information characterizing the unencrypted second (See Van Oorschot Fig. 4 “X-fields” and Col. 6 Lines 45-47); a linker for linking said information and said basic value using the corresponding linkage as it has been used when encrypting to obtain said payload data key (See Van Oorschot Fig. 4 "Unlevelling" and "X-fields" and Col. 4 Lines 54-56); and a second decryptor decrypting the first section containing the encrypted payload data using said payload data key and said payload data encryption algorithm used when encrypting (See Van Oorschot Fig. 4 "symmetric decryption" and "message"), but Van Oorschot failed to disclose that the second section included audio data, video data, a combination of audio data and video data, text data, or binary data forming an executable program, or the X-fields containing audio data, video data, a combination of audio data and video data, text data or binary data forming an executable program.

Nardone teaches that movie data needs to be protected from being copied and that this is generally done through encrypting the movie data (See Nardone Col. 1 Lines 22-37), and further that in order to save on processing cost, only portions of the movie data should be encrypted (See Nardone Col. 1 Summary of the Invention).

Yatsukawa teaches that by providing a public key in a digital certificate, the public key can be authenticated by the recipient of the key (See Yatsukawa Col. 5 Paragraph 1).

It would have been obvious to the ordinary person skilled in the art at the time of invention to employ the teachings of Nardone in the encryption system of Van Oorschot by encrypting video data, and further by only encrypting portions of the data. This would have been obvious because the ordinary person skilled in the art would have been motivated to protect movie data and to save on processing cost. It further would have been obvious to the ordinary person skilled in the art at the time of invention to employ the teachings of Yatsukawa in the encryption system of Van Oorschot by providing the public key in a public key certificate in the X-fields. This would have been obvious because the ordinary person skilled in the art would

have been motivated to provide the recipient with means to verify the authenticity of the public key. Furthermore, it was well know that public key certificates contain text data.

Regarding claim 2, Van Oorschot, Nardone and Yatsukawa disclosed that said payload data encryption algorithm is a symmetrical encryption algorithm (See Van Oorschot Fig. 3 "symmetric encryption").

Regarding claim 3, Van Oorschot, Nardone and Yatsukawa disclosed that said invertible logic linkage is self-inverting and includes an XOR- linkage (See Van Oorschot Col. 6 Lines 56-60).

Regarding claim 4, Van Oorschot, Nardone and Yatsukawa disclosed that one key of said two keys being different from each other is the private key of a producer of said payload data stream or the public key of a consumer of said payload data stream (See Van Oorschot Fig. 3 B's high trust public key).

Regarding claim 5, Van Oorschot, Nardone and Yatsukawa disclosed that said part of said payload data stream being processed to deduce said information includes at least a part of said header (See Van Oorschot Fig. 3 "X-Field" and Col. 6 Lines 49-55).

Regarding claim 6 Van Oorschot, Nardone and Yatsukawa disclosed that said step of processing comprises forming a hash sum (See Van Oorschot Col. 6 Lines 49-55).

Regarding claim 7, Van Oorschot, Nardone and Yatsukawa disclosed further comprising the following step: identifying an algorithm being used in said step of processing by an entry into said header (See Van Oorschot Abstract Lines 14-16).

Regarding claim 14 Van Oorschot, Nardone and Yatsukawa disclosed that said step of processing further comprises the following sub-step: setting said entry for said output value in said header to a defined value and processing said entire header, including said entry set to a defined value (See Van Oorschot Fig. 3 "X-Field" and Col. 6 Lines 49-55).

Regarding Claim 16, Van Oorschot, Nardone and Yatsukawa disclosed the following step: identifying said payload data encryption algorithm by an entry into said header of said payload data stream (See Van Oorschot Abstract Lines 14-16).

Regarding claim 19, Van Oorschot, Nardone and Yatsukawa disclosed that said part being processed to deduce said information is said header (See Van Oorschot Fig. 4 "X-Fields").

Regarding claim 23, Van Oorschot, Nardone and Yatsukawa disclosed that one key having been used when encrypting is the public key of said asymmetrical encryption method, while the other key having been used when decrypting is the private key of said asymmetrical encryption method (See Van Oorschot Fig. 3 "B's high trust public key" and Fig 4 "B's high trust private key").

Regarding claim 24, Van Oorschot, Nardone and Yatsukawa disclosed that said step of processing includes forming a hash sum (See Van Oorschot Col. 6 Lines 49-55 and Fig. 4 "Unlevelling").

Regarding claim 25, Van Oorschot, Nardone and Yatsukawa disclosed that a part of said header having been set to a defined value for said step of processing when encrypting is set to the same defined value for said step of processing when decrypting (See Van Oorschot Fig. 3 "X-fields" and Fig. 4 "X-fields" wherein they must be the same defined value because they were both set by the sender upon sending).

Regarding claim 26, Van Oorschot, Nardone and Yatsukawa disclosed that said part of said header being set to a defined value includes said entry for said output value of said header

(See Van Oorschot Fig. 3 “B’s header field” and Fig. 4 “B’s header field” wherein they must be the same defined value because they were both set by the sender upon sending).

Regarding claim 27, Van Oorschot, Nardone and Yatsukawa disclosed that said step of linking comprises using an XOR-linkage (See Van Oorschot Col. 6 Lines 56-60 and Col. 4 Lines 54-56 and Fig. 4 “Unlevelling”).

Claims 8, 11-12, 18, and 20-21 are rejected under 35 U.S.C. 103(a) as being unpatentable over Van Oorschot, Nardone and Yatsukawa as applied to claims 1 and 17 above, and further in view of Matyas et al. (US Patent Number 5,200,999) hereinafter referred to as Matyas.

Van Oorschot, Nardone and Yatsukawa disclosed a system for sending a message from a sender to a receiver in which the message was encrypted using a key, the key was encrypted, and then the key was sent to the receiver with the encrypted message (See Van Oorschot Abstract and Fig. 3). Van Oorschot further disclosed decrypting the key, and using the key to decrypt the message at the receiver (See Van Oorschot Abstract and Fig. 4). However, Van Oorschot, Nardone and Yatsukawa failed to disclose sending license data along with the key and message.

Matyas teaches that when sending a key, in order to authenticate the use of the key, and the validity of the key, certain data (License data) should be placed in the header along with the key. This data includes key type, key usage data (for history purposes), algorithm identifier, algorithm-specific data, key start date/time, key expiration data/time, device identifier, user identifier, key identifier, logical device identifier, and user-defined data (See Matyas Col. 13 Line 66 – Col. 14 Lines 60). Matyas further teaches that this information should be verified prior to use of the key (See Matyas Col. 100).

It would have been obvious to the ordinary person skilled in the art at the time of invention to employ the teachings of Matyas in the key and message sending system and method of Van Oorschot, Nardone and Yatsukawa by placing the license information, taught by Matyas, in the header of the message and checking this information prior to allowing the key and message to be decrypted. This would have been obvious because the ordinary person skilled in the art would have been motivated to protect the interests of the sender of the message and to ensure the security of the message.

Claim 9 is rejected under 35 U.S.C. 103(a) as being unpatentable over the combination of Van Oorschot, Nardone, Yatsukawa and Matyas as applied to claim 8 above, and further in view of Klemba et al. (US Patent Number 5,710,814) hereinafter referred to as Klemba.

Van Oorschot, Nardone, Yatsukawa and Matyas disclosed sending license data for controlling the usage of a key and message, including usage history (See rejection of claim 8 above), but failed to disclose the data including how often the message could be decrypted.

Klemba teaches that license data can be used to control the number of uses of a cryptographic function (See Klemba Col. 14 Lines 14-19).

It would have been obvious to the ordinary person skilled in the art at the time of invention to employ the teachings of Klemba in the messaging system and method of Van Oorschot, Nardone, Yatsukawa and Matyas by using the license information to limit the number of times the message could be decrypted. This would have been obvious because the ordinary person skilled in the art would have been motivated to protect the interests of the sender of the message as well as to protect the message against compromise.

Claim 10 is rejected under 35 U.S.C. 103(a) as being unpatentable over the combination of Van Oorschot, Nardone, Yatsukawa and Matyas as applied to claim 8 above, and further in view of Edenson et al. (Us Patent Number 6,198,875) hereinafter referred to as Edenson.

Van Oorschot, Nardone, Yatsukawa and Matyas disclosed sending license data for controlling the usage of a key and message, including usage history (See rejection of claim 8 above), but failed to disclose the data including how often the message could be copied and how often it had already been copied.

Edenson teaches that license information can include how many copies of licensed data can be made (See Edenson Col. 4 Paragraph 2).

It would have been obvious to the ordinary person skilled in the art at the time of invention to employ the teachings of Edenson in the messaging system of Van Oorschot, Nardone, Yatsukawa and Matyas by including information regarding the number of allowed copies of the message that are permitted. This would have been obvious because the ordinary person skilled in the art would have been motivated to protect the interests of the message sender, and to protect the message itself from unauthorized distribution. Further, it would have been necessary to also keep track of the number of copies already made in order to enforce the copy limit.

Claim 13 is rejected under 35 U.S.C. 103(a) as being unpatentable over the combination of Van Oorschot, Nardone, Yatsukawa and Matyas as applied to claim 8 above, and further in view of Schneier ("Applied Cryptography Second Edition").

Van Oorschot, Nardone, Yatsukawa and Matyas disclosed sending license data for controlling the usage of a key and message, including usage history (See rejection of claim 8 above), but failed to disclose including the license in the hash function.

Schneier teaches that hashes are used to authenticate the data being hashed upon receipt of the data in order to detect any unauthorized changes to the data (See Schneier Pages 30-31 Section 2.4).

It would have been obvious to the ordinary person skilled in the art at the time of invention to employ the teachings of Schneier in the messaging system of Van Oorschot, Nardone, Yatsukawa and Matyas by hashing the License data along with the X-fields. This would have been obvious because the ordinary person skilled in the art would have been motivated to protect against undetected changes to the license data sent with the message.

Claim 15 is rejected under 35 U.S.C. 103(a) as being unpatentable over Van Oorschot, Nardone, and Yatsukawa as applied to claim 1 above, and further in view of Roediger (US Patent Number 4,899,333).

Van Oorschot, Nardone, and Yatsukawa disclosed sending a message from a sender to a receiver, including a header and a hash of the header (See Van Oorschot Col. 6), but Van Oorschot failed to disclose including a sender identifier and a receiver identifier in the header, or in the hash.

Roediger teaches that packet headers contain a source address (sender identifier) and a destination address (recipient identifier) and that a checksum should include these fields in order to ensure that the fields are not corrupted (See Roediger Col. 37 Lines 53-63).

It would have been obvious to the ordinary person skilled in the art at the time of invention to employ the teachings of Roediger in the messaging system of Van Oorschot, Nardone, and Yatsukawa by including source and destination addresses in the header and including these in the hash. This would have been obvious because the ordinary person skilled in the art would have been motivated to provide means for routing the message from the sender to the receiver and allowing the receiver to verify that it was the intended receiver of the message.

Claim 22 is rejected under 35 U.S.C. 103(a) as being unpatentable over Van Oorschot, Nardone, and Yatsukawa as applied to claim 17 above, and further in view of Schneier.

Van Oorschot, Nardone, and Yatsukawa disclosed using a public key of the receiver for encryption (See rejection of claim 23 above) but failed to disclose using a private key of an asymmetrical key pair for encryption.

Schneier teaches that by encrypting data using a senders private key, the receiver can use the senders public key to authenticate the sender of the data (See Schneier Pages 53-54).

It would have been obvious to employ the teachings of Schneier in the messaging system of Van Oorschot, Nardone, and Yatsukawa by encrypting the leveled key with the private key of the sender and decrypting it with the public key of the sender. This would have been obvious because the ordinary person skilled in the art would have been motivated to provide sender authentication at the receiver.

Claim 30 is rejected under 35 U.S.C. 103(a) as being unpatentable over Van Oorschot, Nardone, and Yatsukawa as applied to claims 28 and 29 above, and further in view of Kane et al. (US Patent Number 5,315,635) hereinafter referred to as Kane.

Van Oorschot, Nardone, and Yatsukawa disclosed sending messages from a sender to a receiver (See Van Oorschot Abstract), but failed to disclose the sending being from a personal computer to a personal computer.

Kane teaches that messages can be sent between personal computers (See Kane Col. 1 Lines 45-51).

It would have been obvious to the ordinary person skilled in the art at the time of invention to employ the teachings of Kane in the messaging system of Van Oorschot, Nardone, and Yatsukawa by sending the encrypted messages from a sending personal computer to receiving personal computer. This would have been obvious because the ordinary person skilled in the art would have been motivated to protect messages sent between two personal computers.

Conclusion

Claims 1-30 have been rejected and claim 31 has been cancelled.

Any inquiry concerning this communication or earlier communications from the examiner should be directed to Matthew T. Henning whose telephone number is (571) 272-3790. The examiner can normally be reached on M-F 8-4.

If attempts to reach the examiner by telephone are unsuccessful, the examiner's supervisor, Ayaz Sheikh can be reached on (571) 272-3795. The fax phone number for the organization where this application or proceeding is assigned is 571-273-8300.

Information regarding the status of an application may be obtained from the Patent Application Information Retrieval (PAIR) system. Status information for published applications may be obtained from either Private PAIR or Public PAIR. Status information for unpublished applications is available through Private PAIR only. For more information about the PAIR system, see http://pair-direct.uspto.gov. Should you have questions on access to the Private PAIR system, contact the Electronic Business Center (EBC) at 866-217-9197 (toll-free). If you would like assistance from a USPTO Customer Service Representative or access to the automated information system, call 800-786-9199 (IN USA OR CANADA) or 571-272-1000.

/Matthew Henning/
Assistant Examiner
Art Unit 2131
6/22/2007

SYED A. ZIA
PRIMARY EXAMINER
06/23/2007